

LECTURE 14

**Theorem 1.** (Greendlinger's Lemma) Let  $D \rightarrow X$  be a  $C(6)$   $[C(4) - T(4)]$  disc diagram. Then either

- 1)  $D$  is trivial
- 2)  $D$  is a single 2-cell
- 3)  $[C(6)]$   $D$  contain a "count" of 6  $i$ -shells/spurs, where  
spurs, 0-shells, 1-shells count as 3,  
2-shells count as 2,  
3-shells count as 1.
- 3)  $[C(4) - T(4)]$   $D$  contain a "count" of 4  $i$ -shells/spurs, where  
spurs, 0-shells, 1-shells count as 2,  
2-shells count as 1.

(figure: examples)

PROOF of  $[C(6)]$ . Assume  $D$  is not a single 2-cell or 0-cell. Ignore all valence 2 0-cells. Assign angles as follows:

(figure) totally internal angles  $\frac{2\pi}{3}$ ,

(figure) half internal angles  $\frac{\pi}{2}$ ,

(figure) totally external angles 0.

Now  $D$  is an angled 2-complex.

- 1) Each internal 0-cell has curvature  $\leq 0$ .
- 2) Each external 0-cell has curvature  $\leq 0$  except endpoint of a spur has curvature  $= \pi$ .

Indeed, if  $v \in \partial D$  but  $\text{valence}(v) \neq 1$ , then the link of  $v$  is a forest with at least 3 vertices. Therefore

either  $\text{link}(v)$  has at least 2 components, so  $\chi(\text{link}(v)) \geq 2$ ,

or  $\text{link}(v)$  is connected and  $\sum \angle \geq \pi$  (figures of 3 possibilities).

So 2-cells must make positive curvature. Each internal 2-cell has non-positive curvature (sum of angles  $\geq \frac{2\pi}{3} \cdot 6$ ). Each 2-cell with  $\geq 2$  boundary arcs has curvature  $\leq 0$ , since has  $\geq 4$  angles of value  $\frac{\pi}{2}$ .

EXERCISE. Check this.

(figures: possible cases)

Remaining 2-cells have exactly 1 boundary edge. Curvature of a 2-cell is  $\kappa(f) = \sum \angle - (|f| - 2)\pi$ . Sum of angles of an  $i$ -shell ( $i \geq 1$ ) is  $\sum \angle = \frac{\pi}{2} + \frac{\pi}{2} + (i - 1)\frac{2\pi}{3}$ . Then  $\kappa(f) = \pi - \frac{\pi}{3}(i - 1)$

$$\kappa(1\text{-shell}) = \pi$$

$$\kappa(2\text{-shell}) = \frac{2\pi}{3}$$

$$\kappa(3\text{-shell}) = \frac{\pi}{3}$$

$$\kappa(4\text{-shell}) = 0$$

(By the way: you may have as many 4-shells as you like.)

$$2\pi\chi(D) = \sum \kappa(v) + \sum \kappa(f)$$

$$\begin{aligned} 2\pi &\leq 0 + \pi \cdot \#_S + \pi \cdot \#_0 + \pi \cdot \#_1 + \frac{2\pi}{3} \cdot \#_2 + \frac{\pi}{3} \cdot \#_3 \\ 6 &\leq 0 + 3 \cdot \#_S + 3 \cdot \#_0 + 3 \cdot \#_1 + 2 \cdot \#_2 + 1 \cdot \#_3 \end{aligned}$$

PROOF SKETCH for  $C(4) - T(4)$  case. Similar argument except use  $\frac{\pi}{2}$  angles on internal and half-internal corners, 0 angles on totally external corners (ends up  $\kappa(f) = (3 - i)\frac{\pi}{2}$ ).  $\square$

Why no lemma for  $C(3) - T(6)$ ? (figure: hexagon) Zillion 2-shells, no 1-shells, no 0-shells. Although, there are positive features. (figure: rhombus).

EXERCISE. Give example of infinitely many  $(6, 3)$  complexes.

**Rips's Short Exact Sequence.** A construction producing interesting subgroups of  $C'(\frac{1}{6})$  groups.

Let  $Q$  be a finitely presented group  $\langle \bar{g}_1, \dots, \bar{g}_s | \bar{R}_1, \dots, \bar{R}_t \rangle$ . There exist a short exact sequence

$$1 \rightarrow N \rightarrow G \rightarrow Q \rightarrow 1$$

s.t.  $G$  is a finitely presented  $C'(\frac{1}{6})$  group,  $N$  is finitely generated.

$$1 \rightarrow \langle x, y \rangle \rightarrow \langle x, y, g_1, \dots, g_s | R_1, \dots, R_t \rangle \rightarrow \langle \bar{g}_1, \dots, \bar{g}_s | \bar{R}_1, \dots, \bar{R}_t \rangle \rightarrow 1$$

Rips's construction can be used to produce  $C'(\frac{1}{6})$  group  $G$  with interesting subgroups from a group  $Q$  with interesting properties.

EXAMPLE. If  $Q$  has undecidable word problem then  $G$  has undecidable membership problem in  $N$ . Induced from a word  $\bar{w}$  in generators of  $Q$ ,  $(\bar{w} = 1) \Leftrightarrow (w \in N)$ .

EXAMPLE. If  $Q$  has a finitely generated subgroup  $\bar{H}$  that is finitely generated but not finitely presented then its preimage  $H$  in  $G$  is finitely generated but not finitely presented.

Indeed, if  $\bar{H} = \langle \bar{w}_1, \dots, \bar{w}_n \rangle$  then  $H = \langle w_1, \dots, w_n, x, y \rangle$ . If  $H$  is finitely presented then  $H = \langle h_1, \dots, h_k | U_1, \dots, U_m \rangle$ , so  $\bar{H} = \langle h_1, \dots, h_k | U_1, \dots, U_m, X, Y \rangle$ .