

# Computation of operators in wavelet coordinates

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# Overview

- Linear operator equation  $Lu = g$  with  $L : \mathcal{H} \rightarrow \mathcal{H}'$
- Riesz basis  $\Psi = \{\psi_\lambda\}$  of  $\mathcal{H}$ , e.g.  $u = \sum_\lambda \mathbf{u}_\lambda \psi_\lambda$
- Infinite dimensional matrix-vector system  $\mathbf{L}\mathbf{u} = \mathbf{g}$ , with  $\mathbf{u} = (\mathbf{u}_\lambda)_\lambda$  and  $\mathbf{L} : \ell_2 \rightarrow \ell_2$
- Convergent iterations such as  $\mathbf{u}^{(i+1)} = \mathbf{u}^{(i)} + \alpha[\mathbf{g} - \mathbf{L}\mathbf{u}^{(i)}]$
- We **can** approximate  $\mathbf{L}\mathbf{u}^{(i)}$  by a finitely supported vector
- How cheap can we **compute** this approximation?
- The answer will depend on  $L$  and  $\Psi$

# Linear operator equations

- Let  $\Omega$  be an  $n$ -dimensional domain or smooth manifold
- $H^t \subset H^t(\Omega)$  be a subspace, and  $H^{-t}$  be its dual space
- Consider the problem of finding  $u$  from

$$Lu = g$$

- where  $L : H^t \rightarrow H^{-t}$  is a self-adjoint elliptic operator of order  $2t$
- and  $g \in H^{-t}$  is a linear functional

# Differential operators

- Partial differential operators of order  $2t$

$$\langle v, Lu \rangle = \sum_{|\alpha|, |\beta| \leq t} \langle \partial^\alpha v, a_{\alpha\beta} \partial^\beta u \rangle,$$

- Example: The reaction-diffusion equation ( $t = 1$ )

$$\langle v, Lu \rangle = \int_{\Omega} \nabla v \cdot \nabla u + \kappa^2 vu,$$

# Singular integral operators

- Boundary integral operators

$$[Lu](x) = \int_{\Omega} K(x, y)u(y)d\Omega_y$$

with the kernel  $K(x, y)$  singular at  $x = y$

- Example: The single layer operator for the Laplace BVP in 3-d domain ( $t = -\frac{1}{2}$ )

$$K(x, y) = \frac{1}{4\pi|x - y|}$$

# Multiresolution analysis

- $\mathcal{S}_0 \subset \mathcal{S}_1 \subset \dots \subset H^t$  and  $\tilde{\mathcal{S}}_0 \subset \tilde{\mathcal{S}}_1 \subset \dots \subset H^{-t}$
- $\dim \mathcal{S}_j, \dim \tilde{\mathcal{S}}_j = \mathcal{O}(2^{jn})$  (dyadic refinements)
- $\mathcal{S}_j$  contains all piecewise pols of degree  $d - 1$
- $\tilde{\mathcal{S}}_j$  contains all piecewise pols of degree  $\tilde{d} - 1$
- $\mathcal{S}_j$  is globally  $C^r$ -smooth

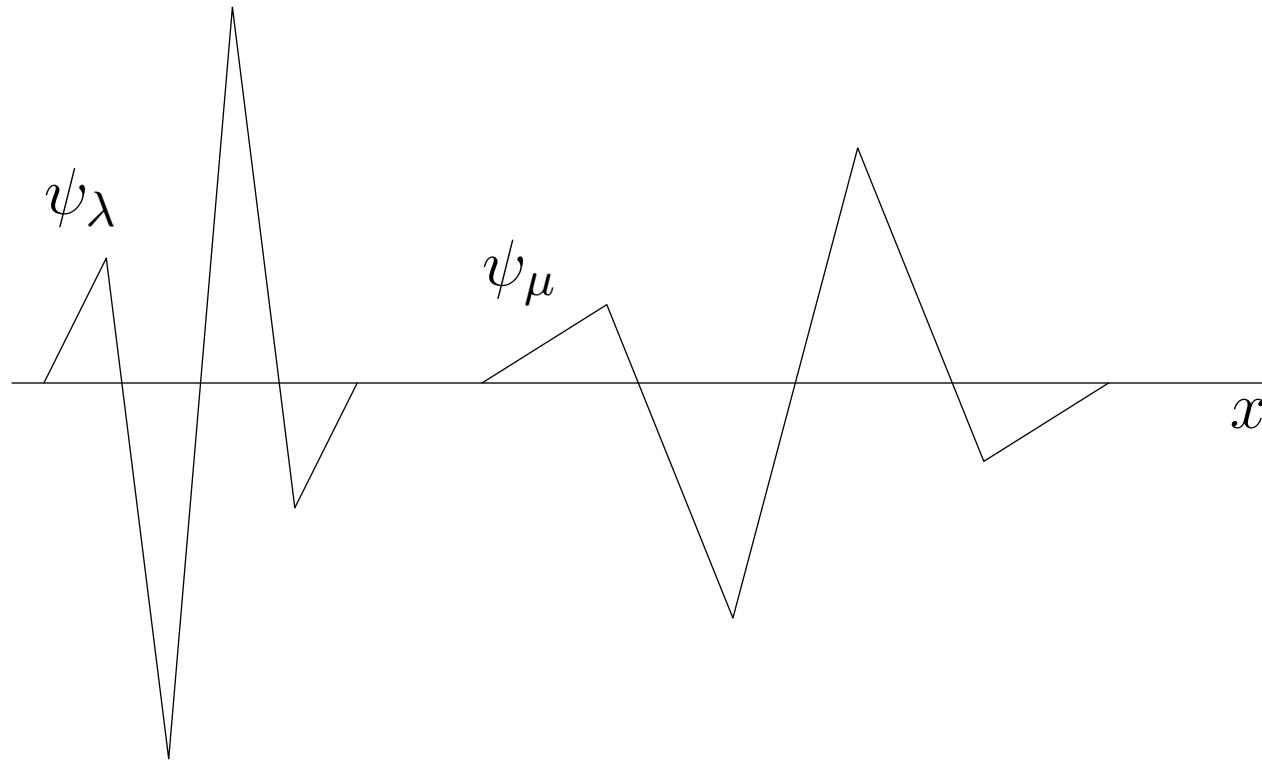
# Wavelet bases

- $\Psi = \{\psi_\lambda : \lambda \in \Lambda\}$  is a Riesz basis for  $H^t$ 
  - each  $v \in H^t$  has a unique expansion

$$v = \sum_{\lambda \in \Lambda} \mathbf{v}_\lambda \psi_\lambda \quad \text{s.t.} \quad c\|\mathbf{v}\| \leq \|v\|_{H^t} \leq C\|\mathbf{v}\|$$

- For every index  $\lambda \in \Lambda$ , there is a number  $|\lambda| \in \mathbb{N}_0$  called the **level** of  $\psi_\lambda$
- $\text{span}\{\psi_\lambda : |\lambda| \leq j\} = \mathcal{S}_j$
- $\langle \psi_\lambda, v \rangle = 0$  for any  $v \in \tilde{\mathcal{S}}_{|\lambda|-1}$
- $\text{diam}(\text{supp } \psi_\lambda) = \mathcal{O}(2^{-|\lambda|})$

# Typical wavelets



- $\psi_\lambda$  is a piecewise polynomial of degree  $d - 1$
- $\int x^k \psi_\lambda(x) dx = 0$  for  $k < \tilde{d}$  ( $\tilde{d}$  vanishing moments)

# Galerkin methods

- Wavelet basis  $\Psi_j := \{\psi_\lambda : |\lambda| \leq j\}$  of  $\mathcal{S}_j$
- Stiffness  $\mathbf{L}_{(j)} = \langle L\psi_\lambda, \psi_\mu \rangle_{|\lambda|, |\mu| \leq j}$
- load  $\mathbf{g}_{(j)} = \langle g, \psi_\lambda \rangle_{|\lambda| \leq j}$
- Linear equation in  $\mathbb{R}^{N_j}$  ( $N_j := \dim \mathcal{S}_j$ )

$$\mathbf{L}_{(j)} \mathbf{u}_{(j)} = \mathbf{g}_{(j)}$$

- $\mathbf{L}_{(j)} : \mathbb{R}^{N_j} \rightarrow \mathbb{R}^{N_j}$  SPD and  $\mathbf{g}_{(j)} \in \mathbb{R}^{N_j}$
- $u_{(j)} = \sum_\lambda [\mathbf{u}_{(j)}]_\lambda \psi_\lambda$  approximates the solution of  $Lu = g$

# Galerkin approximation

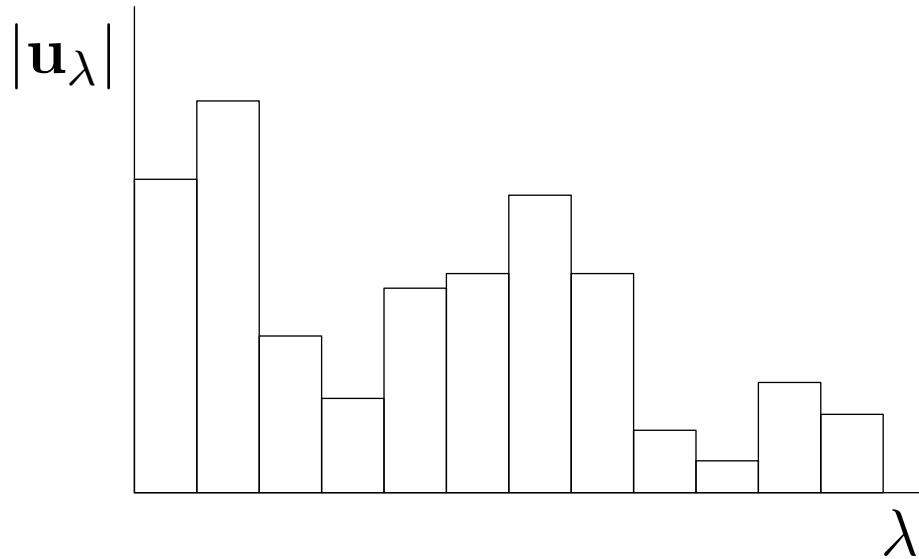
- If  $u \in H^s$  for some  $s \in [t, d]$

$$\varepsilon_{(j)} := \|u_{(j)} - u\|_{H^t} \leq \mathcal{O}(2^{-j(s-t)})$$

- $N_j = \dim \mathcal{S}_j = \mathcal{O}(2^{jn})$
- $\varepsilon_{(j)} \leq \mathcal{O}(N_j^{-\frac{s-t}{n}})$
- Solve  $\mathbf{L}_{(j)} \mathbf{u}_{(j)} = \mathbf{g}_{(j)}$  with CG  $\leadsto$  complexity  $\mathcal{O}(N_j)$
- Similar estimates for FEM
- Better convergence? Adaptive methods?

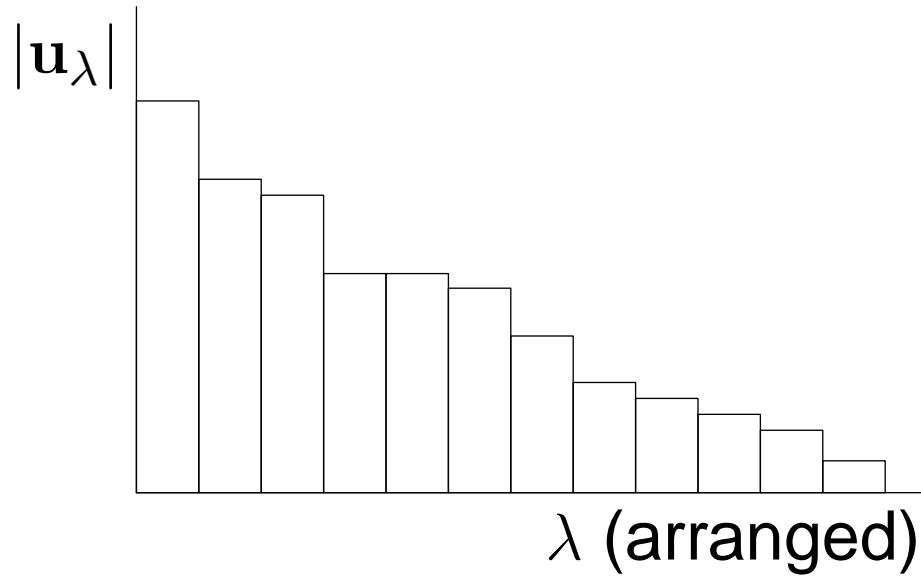
# Nonlinear approximation

- Given  $\mathbf{u} = (\mathbf{u}_\lambda)_\lambda \in \ell_2$
- Approximate  $\mathbf{u}$  using  $N$  coeffs



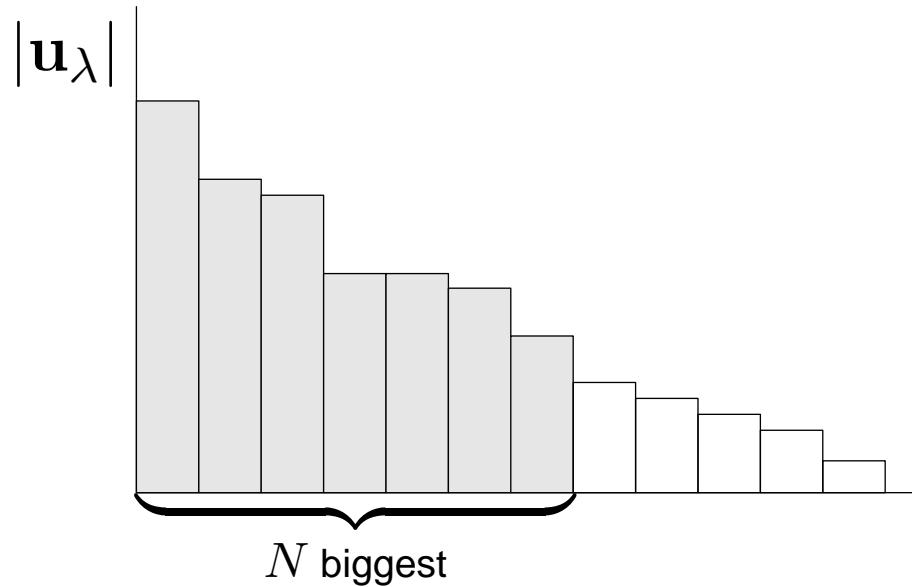
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# Nonlinear approximation

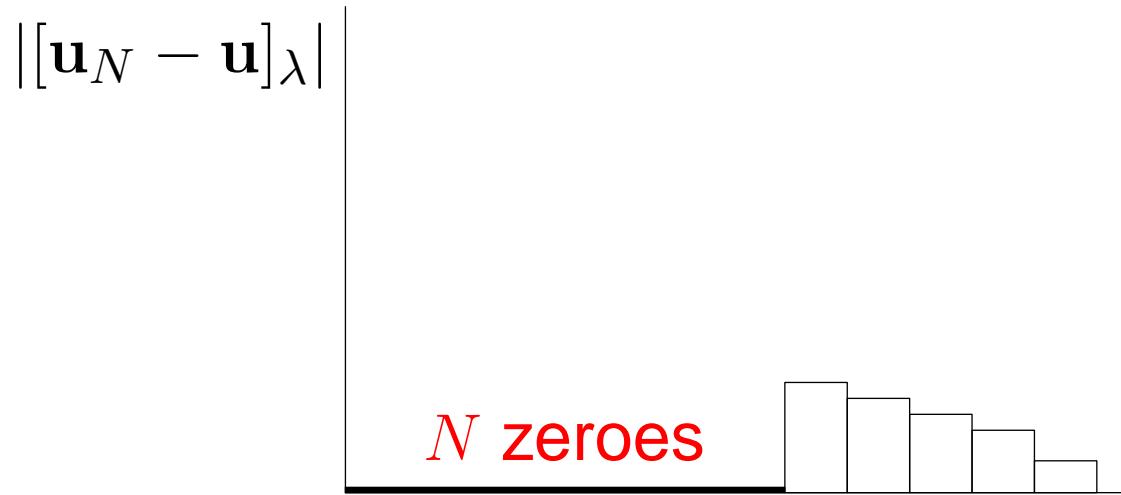
- Given  $\mathbf{u} = (\mathbf{u}_\lambda)_{\lambda} \in \ell_2$
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- $\mathbf{u}_N$  best approximation of  $\mathbf{u}$  with  $\#\text{supp } \mathbf{u}_N \leq N$

# Nonlinear approximation

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- $\mathbf{u}_N$  best approximation of  $\mathbf{u}$  with  $\#\text{supp } \mathbf{u}_N \leq N$

# Nonlinear vs. linear approximation

- If  $u \in B_{\tau,\tau}^s$  with  $\frac{1}{\tau} = \frac{1}{2} + \frac{s-t}{n}$  for some  $s < d$

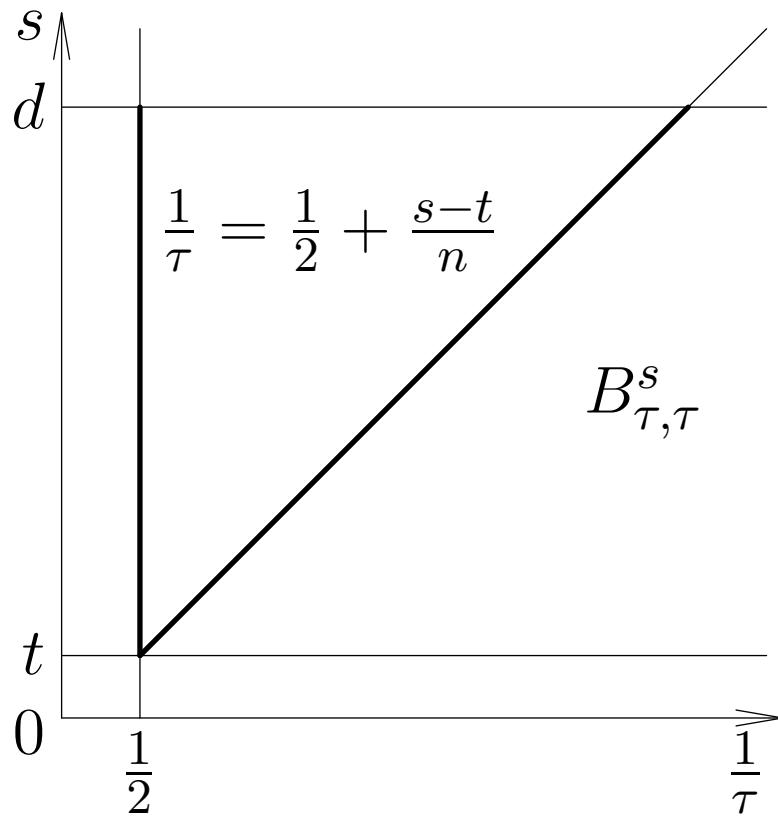
$$\varepsilon_N = \|\mathbf{u}_N - \mathbf{u}\| \leq \mathcal{O}(N^{-\frac{s-t}{n}})$$

- If  $u \in H^s$  for some  $s \leq d$ , uniform refinement

$$\varepsilon_{(j)} = \|\mathbf{u}_{(j)} - \mathbf{u}\| \leq \mathcal{O}(N_j^{-\frac{s-t}{n}})$$

- $B_{\tau,\tau}^s$  is bigger than  $H^s$

# Besov vs. Sobolev regularity



- [Dahlke, DeVore]:  $u \in B_{\tau,\tau}^d$  with  $\frac{1}{\tau} = \frac{1}{2} + \frac{d-t}{n}$  "often"

# Equivalent problem in $\ell_2$

[Cohen, Dahmen, DeVore '02]

- Wavelet basis  $\Psi = \{\psi_\lambda : \lambda \in \Lambda\}$
- **Stiffness**  $\mathbf{L} = \langle L\psi_\lambda, \psi_\mu \rangle_{\lambda, \mu}$  and **load**  $\mathbf{g} = \langle g, \psi_\lambda \rangle_\lambda$
- Linear equation in  $\ell_2(\Lambda)$

$$\mathbf{Lu} = \mathbf{g}$$

- $\mathbf{L} : \ell_2(\Lambda) \rightarrow \ell_2(\Lambda)$  SPD and  $\mathbf{g} \in \ell_2(\Lambda)$
- $u = \sum_\lambda \mathbf{u}_\lambda \psi_\lambda$  is **the solution** of  $Lu = g$

# Richardson iterations in $\ell_2$

- $\mathbf{u}^{(0)} = \mathbf{0}$
  - $\mathbf{u}^{(i+1)} = \mathbf{u}^{(i)} + \alpha[\mathbf{g} - \mathbf{L}\mathbf{u}^{(i)}] \quad i = 0, 1, \dots$
  - $\mathbf{g}$  and  $\mathbf{L}\mathbf{u}^{(i)}$  are **infinitely** supported
  - Approximate them by **finitely** supported sequences
  - Algorithm **SOLVE**[ $N, \mathbf{L}, \mathbf{g}$ ]  $\rightarrow \mathbf{u}_{[N]}$  ( $N$  operations)
  - $\#\text{supp } \mathbf{u}_{[N]} \leq \mathcal{O}(N)$  and
- $$\varepsilon_{[N]} = \|\mathbf{u}_{[N]} - \mathbf{u}\| \rightarrow 0 \quad \text{as } N \rightarrow \infty$$
- $\varepsilon_{[N]}$  speed of convergence?

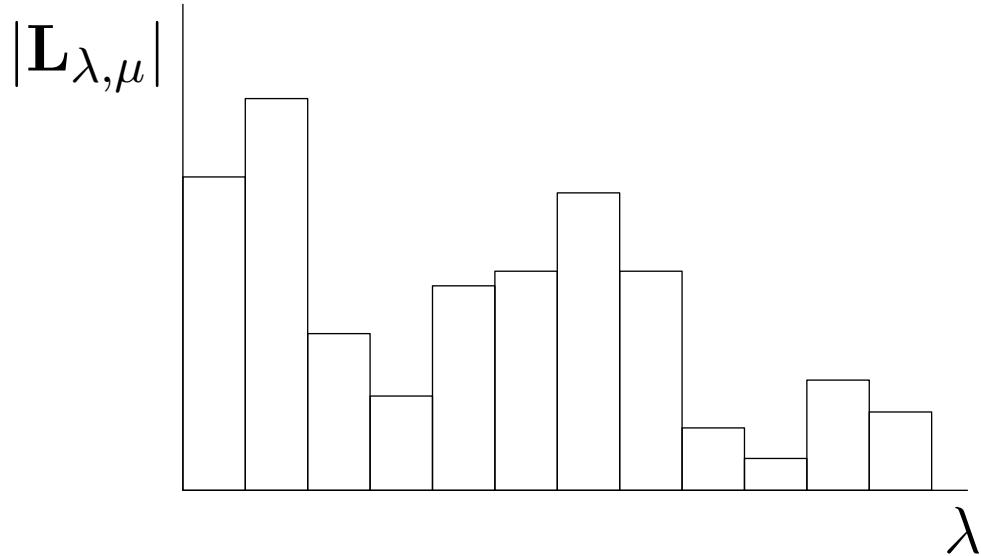
# Complexity of SOLVE

- Matrix  $\mathbf{L}$  is called  $q^*$ -computable, when for each  $N$  one can construct an infinite matrix  $\mathbf{L}_N$  s.t.
  - for any  $q < q^*$ ,  $\|\mathbf{L}_N - \mathbf{L}\| \leq \mathcal{O}(N^{-q})$
  - having in each column  $\mathcal{O}(N)$  non-zero entries
  - whose computation takes  $\mathcal{O}(N)$  operations
- [CDD'02]: Suppose that
  - $\|\mathbf{u}_N - \mathbf{u}\| \leq \mathcal{O}(N^{-s}) \quad [s < \frac{d-t}{n}]$
  - $\mathbf{L}$  is  $q^*$ -computable with  $q^* > s$then for suitable  $\mathbf{g}$ ,  $\mathbf{u}_{[N]} = \text{SOLVE}[N, \mathbf{L}, \mathbf{g}]$  satisfies

$$\|\mathbf{u}_{[N]} - \mathbf{u}\| \leq \mathcal{O}(N^{-s})$$

# Computability

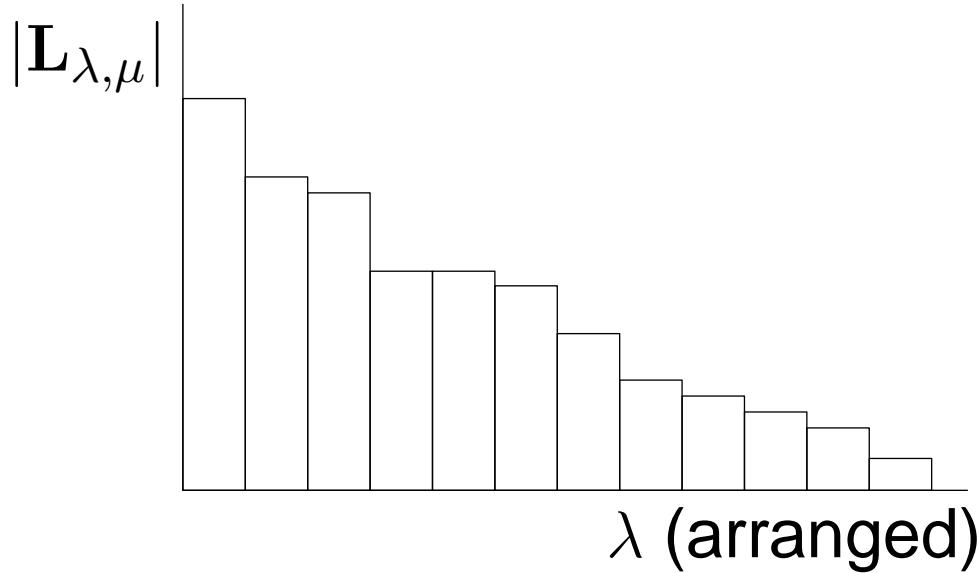
- $[\mathbf{L}_{\lambda,\mu}]_{\lambda \in \Lambda}$  –  $\mu$ -th column



- Approximate by  $N$  entries?

# Computability

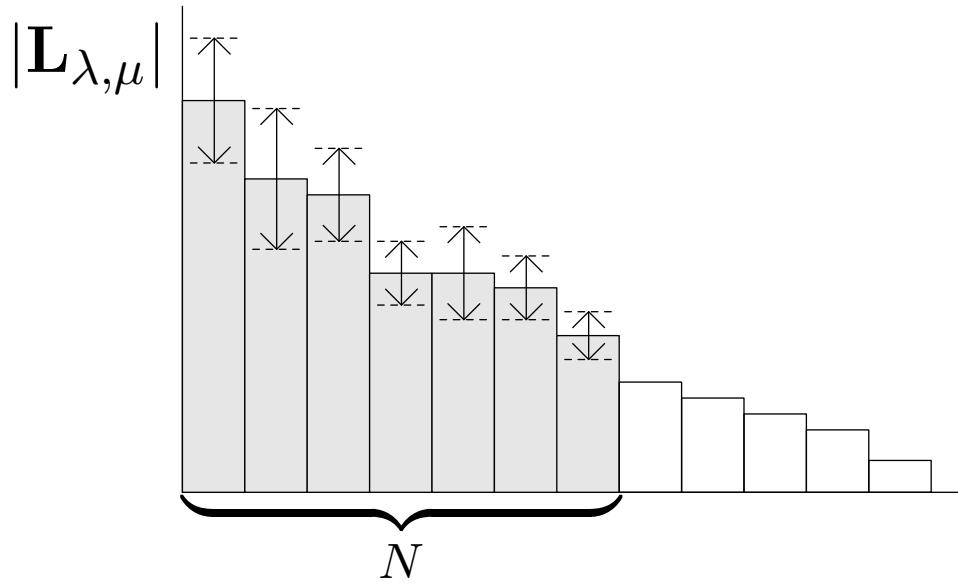
- $[\mathbf{L}_{\lambda,\mu}]_{\lambda \in \Lambda}$  –  $\mu$ -th column arranged by modulus



- $N$  biggest entries?

# Computability

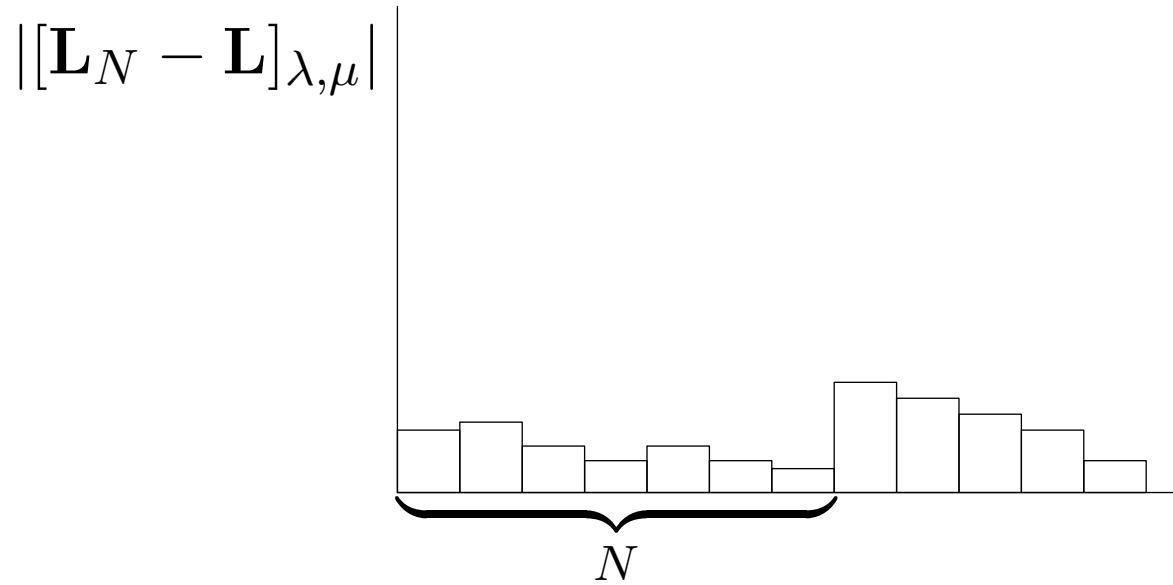
- $[\mathbf{L}_{\lambda,\mu}]_{\lambda \in \Lambda} - \mu\text{-th column}$



- Compute the  $N$  biggest entries

# Computability

- The  $\mu$ -th column of the difference



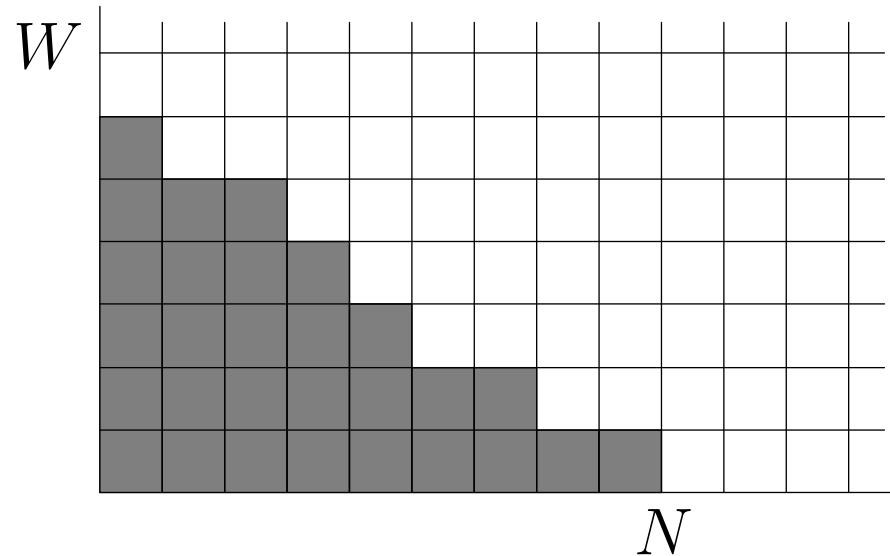
- Need to locate the biggest entries **a priori**

# Compressibility

- $\mathbf{L}$  is called  $q^*$ -compressible, when  $\mathbf{L}$  is  $q^*$ -computable assuming each entry of  $\mathbf{L}$  is available at unit cost
  - [CDD'01], [Stevenson '04]: Suppose
    - $\{\psi_\lambda\}$  are piecewise polynomial wavelets that
      - are sufficiently smooth and
      - have sufficiently many vanishing moments
    - $\mathbf{L}$  is either differential or singular integral operator
- then  $\mathbf{L}$  is  $q^*$ -compressible for some  $q^* \geq \frac{d-t}{n} \quad (> s)$

# Computability

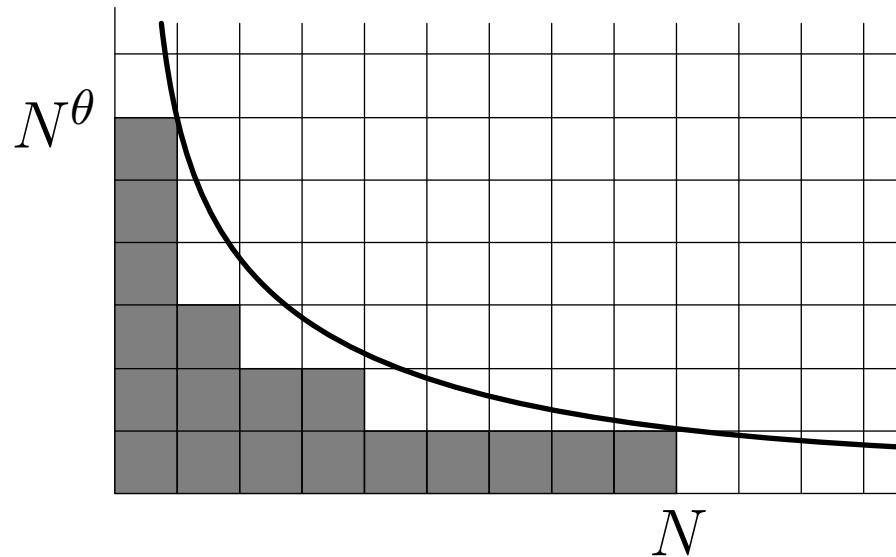
- Distribute computational works over the entries



- Require: shaded area =  $\mathcal{O}(N)$

# Computability

- Distribute computational works over the entries



- $p(x) = N^\theta x^{-\varrho} \sim$  when  $\theta \leq \varrho < 1$

$$\int_0^N N^\theta x^{-\varrho} dx = \frac{1}{1-\varrho} N^{1+\theta-\varrho} = \mathcal{O}(N)$$

# Computability

[T.G., Stevenson '04], [T.G.'04]: Suppose

- $\{\psi_\lambda\}$  are piecewise polynomial wavelets that
  - are sufficiently smooth and
  - have sufficiently many vanishing moments
- $L$  is either differential or singular integral operator

then  $L$  is  $q^*$ -computable for some  $q^* \geq \frac{d-t}{n}$  ( $> s$ )

So the adaptive wavelet method has the optimal convergence rate and optimal computational complexity

# References

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