## Math 595: Descriptive Set Theory HOMEWORK 3 Due: Week of Oct 24

- 1. Show that for any Polish space X there is a continuous open surjection  $g : \mathbb{N}^{\mathbb{N}} \to X$  by constructing a sequence  $(U_s)_{s \in \mathbb{N}^{<\infty}}$  of open subsets of X such that
  - (i)  $U_{\emptyset} = X$

(ii) 
$$\overline{U}_{s^{\frown}i} \subseteq U_s$$

- (iii)  $U_s = \bigcup_i U_{s^{\frown}i}$
- (iv) diam $(U_s) < 2^{-|s|}$ .

CAUTION: We don't require  $U_{s^{i}} \cap U_{s^{j}} = \emptyset$  for  $i \neq j$  (which makes your life easy), so the associated map g may not be injective.

- **2.** The following steps outline a proof of the Baire category theorem for locally compact Hausdorff spaces.
  - (i) [*Optional*] Show that compact Hausdorff spaces are normal.
  - (ii) [*Optional*] Using part (i), prove that in locally compact<sup>1</sup> Hausdorff space X, for every nonempty open set U and every point  $x \in U$ , there is a nonempty precompact<sup>2</sup> open  $V \ni x$  with  $\overline{V} \subseteq U$ .
  - (iii) Prove that locally compact Hausdorff spaces are Baire.
- **3.** For topological spaces *X*, *Y*, a continuous map  $f : X \rightarrow Y$  is called *category preserving* if *f*-preimages of meager sets are meager.
  - (a) Show that any continuous open map  $f : X \to Y$  is category preserving (in fact, f-preimages of nowhere dense are nowhere dense). In particular, projections are category preserving.
  - (b) For topological spaces *X*, *Y*, if *X* is Baire, then, for a continuous map  $f : X \to Y$ , the following are equivalent:
    - (1) *f* is category preserving.
    - (2) *f*-preimages of nowhere dense sets are nowhere dense.
    - (3) *f*-images of open sets are somewhere dense.
    - (4) *f*-preimages of dense open sets are dense.

<sup>&</sup>lt;sup>1</sup>A topological space is said to be **locally compact** if every point has a neighborhood basis that consists of precompact<sup>2</sup> open sets.

<sup>&</sup>lt;sup>2</sup>**Precompact** sets are those contained in compact sets. For Hausdorff spaces, this is equivalent to having a compact closure.

- 4. (a) [*Optional*] Give an example of a function  $f : \mathbb{R} \to \mathbb{R}$  that is continuous at every irrational but discontinuous at every rational.
  - (b) Prove that there is no function  $f : \mathbb{R} \to \mathbb{R}$  that is continuous at every rational but discontinuous at every irrational.

HINT: Show that the set of continuity points of any function is  $G_{\delta}$ .

REMARK: This is yet another example of a proof of nonexistence of an object based on the complexity of a set associated with it. I'm a fan.

- **5.** [*Optional, but at least read it.*] Recall that C([0,1]) is a Polish space with the uniform metric. Show that a generic element of C([0,1]) is nowhere differentiable following the outline below.
  - (i) Prove that given  $m \in \mathbb{N}$ , any function  $f \in C([0,1])$  can be approximated (in the uniform metric) by a piecewise linear function  $g \in C([0,1])$ , whose linear pieces (finitely many) have slope  $\pm M$ , for some  $M \ge m$ .
  - (ii) For each  $n \ge 1$ , let  $E_n$  be the set of all functions  $f \in C([0,1])$ , for which there is  $x_0 \in [0,1]$  (depending on f) such that  $|f(x) f(x_0)| \le n|x x_0|$  for all  $x \in [0,1]$ . Show that  $E_n$  is nowhere dense using the fact that if g is as in (1) with m = 2n, then some open neighborhood of g is disjoint from  $E_n$ .
- 6. Let X be a perfect Polish space and show that a generic compact subset of X is perfect; in fact, the set  $\mathcal{K}_p(X)$  of all perfect compact subsets of X is dense  $G_{\delta}$  in  $\mathcal{K}(X)$ . Conclude that X has continuum many perfect compact subsets.

HINT: Let  $\mathcal{U}$  be a countable open basis for X, and for each  $U \in \mathcal{U}$ , let  $\mathcal{K}_U$  denote the set of  $K \in \mathcal{K}(X)$  such that either  $K \cap U = \emptyset$  or  $|K \cap U| \ge 2$ . Prove that each  $\mathcal{K}_U$  is dense  $G_{\delta}$ .

- 7. Let  $\mathcal{G}$  be the so-called *Hamming graph* on  $2^{\mathbb{N}}$ , namely, there is an edge between  $x, y \in 2^{\mathbb{N}}$  exactly when x and y differ by one bit.
  - (a) Prove that  $\mathcal{G}$  is has no odd cycles and hence is bipartite (admits a 2-coloring). Pinpoint the use of AC.
  - (b) Fix a coloring  $c: 2^{\mathbb{N}} \to 2$  of  $\mathcal{G}$  and let  $A_i := c^{-1}(i)$  for  $i \in \{0, 1\}$ . Consider the game where each player plays a finite nonempty binary sequence at each step and a run of the game is the concatenation of those finite sequences, thus an infinite binary sequence. Prove that this game with the payoff set  $A_0$  is not determined by showing that if one of the players had a winning strategy, so would the other one.

HINT: Steal the other player's strategy.

8. A *finite bounded game* on a set *A* is a game similar to infinite games, but the players play at most *n* number of steps before the winner is decided, for some fixed number  $n \ge 1$  (say a million). More formally, the game is a tree  $T \subseteq A^{<n}$ , for some *n*, and the runs of the game are exactly the elements of the set Leaves(*T*) of all leaves of *T*, so the payoff set is a subset  $D \subseteq \text{Leaves}(T)$ . Player I wins the run  $s \in \text{Leaves}(T)$  of the game iff  $s \in D$ . Consequently, Player II wins iff  $s \in \text{Leaves}(T) \setminus D$ . All games that appear in real life are such games, e.g. chess (counting ties as a win for Player II).

Prove the determinacy of finite bounded games.

HINT: Let's write down what it means for Player I to have a winning strategy in this game, assuming for simplicity that *n* is even and that all of the runs of the game are of length exactly *n*:

 $\exists a_1 \forall a_2 \dots \exists a_{n-1} \forall a_n ((a_1, \dots, a_n) \in D).$ 

What happens when you negate this statement?

- **9.** A *finite game* on a set *A* is a game similar to infinite games, but the players play only finitely many steps before the winner is decided. More formally, it is a (possibly infinite) tree  $T \subseteq A^{<\mathbb{N}}$  that has no infinite branches, and the set of runs is Leaves(*T*), so the payoff set is a subset  $D \subseteq \text{Leaves}(T)$ . Player I wins the run  $s \in \text{Leaves}(T)$  of the game iff  $s \in D$ . Consequently, Player II wins iff  $s \in \text{Leaves}(T) \setminus D$ .
  - (a) Prove the determinacy of finite games.

HINT: Call a position  $s \in T$  determined, if from that point on, one of the players has a winning strategy. Thus, no player has a winning strategy in the beginning iff  $\emptyset$  is undetermined. What can you say about extensions of undetermined positions?

- (b) Conclude the determinacy of clopen infinite games. (These are games with runs in  $A^{\mathbb{N}}$  and the payoff set a clopen subset of  $A^{\mathbb{N}}$ .)
- **10.** [*Optional*] In ZF (in particular, don't use AC or  $\neg$ AD), define a game with rules G(T,D) on the set  $A = \mathscr{P}(\mathbb{N}^{\mathbb{N}})$  (i.e. define a pruned tree  $T \subseteq A^{<\mathbb{N}}$  and a set  $D \subseteq A^{\mathbb{N}}$ ), so that ZF+ $\neg$ AD implies that this game is undetermined. In other words, you have to define the tree *T* and the payoff set *D* without using  $\neg$ AD, but then prove that the game G(T,D) is undetermined using  $\neg$ AD.

HINT: Note that besides playing subsets of  $\mathbb{N}^{\mathbb{N}}$ , players can also play natural numbers in the sense that  $\mathbb{N} \hookrightarrow \mathscr{P}(\mathbb{N}^{\mathbb{N}})$  by  $n \mapsto \{(n)_{i \in \mathbb{N}}\}$ .

- **11.** Prove that a topological group *G* is Baire iff *G* is nonmeager.
- **12.** Let *X* be a topological space and  $A \subseteq X$ . Prove:
  - (a) U(A) is regular open, i.e. it is equal to the interior of its closure.
  - (b) If moreover X is a Baire space and A is Baire measurable, then U(A) is the unique regular open set U with A = U.