

MATH 570: MATHEMATICAL LOGIC

POTENTIAL FINAL EXAM PROBLEMS

1. Let $(L, <)$ be a linearly ordered set and let $(A_\ell)_{\ell \in L}$ be an increasing elementary chain of σ -structures, i.e. $A_\ell \leq A_k$ for $\ell \leq k$. Show that $A = \bigcup_{\ell \in L} A_\ell$ is a universe of a σ -structure A and that $A_\ell \leq A$ for each $\ell \in L$.
2. Let M be a σ -structure.

Definition. M is said to be *weakly homogeneous* if for every pair $A \subseteq B$ of finitely generated substructures of M , every embedding (not necessarily elementary) $h : A \hookrightarrow M$ extends to an embedding $\bar{h} : B \hookrightarrow M$.

Definition. For a subset $A \subseteq M$ and a σ -structure K , a function $h : A \rightarrow K$ is called a *partial elementary map* $M \rightarrow_e K$ if for every σ -formula $\varphi(\vec{x})$ and $\vec{a} \in A^{|\vec{x}|}$,

$$M \models \varphi(\vec{a}) \text{ if and only if } K \models \varphi(h(\vec{a})).$$

Definition. For an infinite cardinal κ , M is said to be κ -*homogeneous* if for every $A \subseteq M$ with $|A| < \kappa$ and $a \in M$, every partial elementary map $h : A \rightarrow M$ extends to a partial elementary map $\bar{h} : A \cup \{a\} \rightarrow M$.

Prove that if M is weakly homogeneous, then for every finitely generated substructure $A \subseteq M$, any embedding $A \hookrightarrow M$ is a partial elementary map $M \rightarrow M$. Deduce that weak homogeneity implies \aleph_0 -homogeneity.

3. Give an example of a structure A (in some signature) and a definable binary relation Q in it such that the unary relation

$$P(y) := \Leftrightarrow \exists^\infty x Q(x, y)$$

is not definable in A , where “ $\exists^\infty x$ ” means “for infinitely many x ”. Provide a proof of every statement you use that claims nondefinability of a set.

4. Let $\sigma := (E)$ be the signature of graphs, i.e. E is a binary relation symbol.
 - (a) Write down an explicit axiomatization T for the class of undirected graphs with no loops, whose connected components are *bi-infinite chains*, i.e. acyclic graphs with the degree of each vertex being 2.
 - (b) Show that T is complete.
 - (c) Conclude yet again (for the last time, I promise) that the binary relation R of being in the same connected component is not 0-definable in any disconnected model of T .

HINT: Let $\varphi_R(x, y)$ be a formula defining R in a disconnected model M , so $M \models \exists x \exists y \neg \varphi_R(x, y)$, hence $Z \models \exists x \exists y \neg \varphi_R(x, y)$, where Z is the connected model. Elements $a, b \in Z$ that witness the latter sentence are within finite distance from each other...

- (d) Show that for any $M \models T$ and $a, b \in M$, there is an automorphism g of M with $g(a) = b$.
- (e) For $M \models T$, exactly which subsets of M are 0-definable in M ?

- (f) Finally, prove that T is model-complete, but does not admit q.e. Recall that that $\text{FOIL}(\sigma)$ includes the 0-ary relation symbols for truth and falsehood, so σ not having a constant symbol is not the reason why T doesn't admit q.e.

5. Consider the following function:

$$f(n) = \begin{cases} 1 & \text{if the decimal expression of } \pi \text{ contains } n \text{ consecutive zeroes} \\ 0 & \text{otherwise} \end{cases}.$$

Is it Σ_1^0 ? Recursive? Primitive recursive? You may assume that the function that outputs the n^{th} decimal digit of π is primitive recursive.

6. Give an example of a Σ_1^0 unary relation that is not recursive. Justify your answer.
7. For each of the following, either prove that it holds for all σ_{arithm} -sentences θ or provide an example of a θ for which it fails. Justify your answers.
- (a) $\text{PA} \models \theta \rightarrow \text{Provable}_{\text{PA}}([\theta])$.
- (b) $\text{PA} \models \text{Provable}_{\text{PA}}([\theta]) \rightarrow \theta$.
- (c) If $N \models \neg \text{Provable}_{\text{PA}}([\theta])$ then $\text{PA} \models \neg \text{Provable}_{\text{PA}}([\theta])$.
- (d) If $\text{PA} \models \neg \text{Provable}_{\text{PA}}([\theta])$ then $N \models \neg \text{Provable}_{\text{PA}}([\theta])$.

8. Let $A, B \subseteq \mathbb{N}^k$. Prove:

- (a) Reduction property for Σ_1^0 : If A, B are Σ_1^0 , then there are disjoint Σ_1^0 sets $A^*, B^* \subseteq \mathbb{N}^k$ such that $A^* \subseteq A$, $B^* \subseteq B$ and $A^* \cup B^* = A \cup B$.

HINT: For $\vec{a} \in \mathbb{N}^k$, decide whether to put it in A^* or B^* based on which of A and B claims it first (i.e. has the smaller witness).

- (b) Separation property for Π_1^0 : If A, B are disjoint Π_1^0 sets, then there is a Δ_1^0 (and hence recursive) set $S \subseteq \mathbb{N}^k$ such that $S \supseteq A$ and $S^c \supseteq B$.

9. Let $n, k \in \mathbb{N}$.

- (a) Construct universal sets for $\Sigma_n^0(\mathbb{N}^k)$ and $\Pi_n^0(\mathbb{N}^k)$.

HINT: This is in the notes for $n = 1$ and the rest is by induction.

- (b) Prove that $\Delta_n^0(\mathbb{N}^k)$ does not admit a universal set.

- (c) Deduce that $\Delta_n^0(\mathbb{N}^k) \subsetneq \Sigma_n^0(\mathbb{N}^k) \subsetneq \Delta_{n+1}^0(\mathbb{N}^k)$ and $\Delta_n^0(\mathbb{N}^k) \subsetneq \Pi_n^0(\mathbb{N}^k) \subsetneq \Delta_{n+1}^0(\mathbb{N}^k)$. Make sure to also show the inclusions, not just their strictness.

- (d) Show that $\bigcup_n \Sigma_n^0 = \bigcup_n \Delta_n^0 = \bigcup_n \Pi_n^0$ is precisely the class of all arithmetical sets.

- (e) Conclude Tarski's theorem that $\ulcorner \text{Th}(N) \urcorner$ is not arithmetical, where $N := (\mathbb{N}, 0, S, +, \cdot)$.

10. Prove that for a σ -theory T , the following are equivalent:

- (1) T is model-complete.
- (2) For every model $A \models T$, $T \cup \text{Diag}(A)$ is a complete σ_A -theory.
- (3.a) Every σ -formula $\varphi(\vec{x})$ is equivalent in T to a universal formula.
- (3.b) Every σ -formula $\varphi(\vec{x})$ is equivalent in T to an existential formula.

HINT: For (2) \Rightarrow (3.a), mimic the proof of “diagram-complete \implies q.e.” More precisely, consider the set

$$\Gamma(\vec{x}) := \{\psi : \psi(\vec{x}) \text{ is a universal } \sigma\text{-formula and } T \models \varphi \rightarrow \psi\}$$

and show that $T \cup \Gamma(\vec{x}) \models \varphi$ (the Generalization axiom is involved).

11. Let \mathcal{C} be a class (a set) of σ -structures. Define the *theory* $\text{Th}(\mathcal{C})$ and the *asymptotic theory* $\text{Th}_a(\mathcal{C})$ of \mathcal{C} as follows: for every σ -sentence φ ,

$$\varphi \in \text{Th}(\mathcal{C}) :\Leftrightarrow \forall \mathbf{M} \in \mathcal{C} \ \mathbf{M} \models \varphi,$$

$$\varphi \in \text{Th}_a(\mathcal{C}) :\Leftrightarrow \forall^\infty \mathbf{M} \in \mathcal{C} \ \mathbf{M} \models \varphi,$$

where \forall^∞ means “for all but finitely many”.

- (a) Let \mathcal{C} be an infinite class of finite structures that contains only finitely many structures of cardinality n , for each $n \in \mathbb{N}$. Prove that the models of $\text{Th}_a(\mathcal{C})$ are exactly the infinite models of $\text{Th}(\mathcal{C})$.

HINT: To prove that any infinite model $\mathbf{M} \models \text{Th}(\mathcal{C})$ is a model of $\text{Th}_a(\mathcal{C})$, note that if $\varphi \in \text{Th}_a(\mathcal{C})$, then $\psi \rightarrow \varphi \in \text{Th}_a(\mathcal{C})$ for any ψ . Find a suitable ψ that is true in \mathbf{M} and $\psi \rightarrow \varphi \in \text{Th}(\mathcal{C})$.

- (b) Let $\sigma_{\text{gr}} := (E)$ be the signature of graphs and for each $n \geq 2$, let \mathbf{G}_n be the (un-directed) graph that is a chain (path) of n vertices, i.e. $\mathbf{G}_n := (V_n, E^{\mathbf{G}_n})$, where $V_n := \{1, 2, \dots, n\}$ and $uE^{\mathbf{G}_n}v :\Leftrightarrow |u - v| = 1$, for $u, v \in V_n$. Let $\mathcal{C} := \{\mathbf{G}_n : n \geq 2\}$. Let

$$T := \{\chi, \theta\} \cup \{\varphi_n, \psi_n : n \geq 2\},$$

where

- χ says that E is irreflexive and symmetric.
- θ says that there are exactly two vertices with degree 1 (call them *leaves*) and all other vertices have degree 2.
- φ_n says that there are two leaves whose distance is at least n .
- ψ_n says that there are no cycles of length n .

Prove that $T \subseteq \text{Th}_a(\mathcal{C})$.

- (c) Prove that T is complete.
(d) Conclude that T is an axiomatization of $\text{Th}_a(\mathcal{C})$.
(e) Describe all models of $\text{Th}_a(\mathcal{C})$.