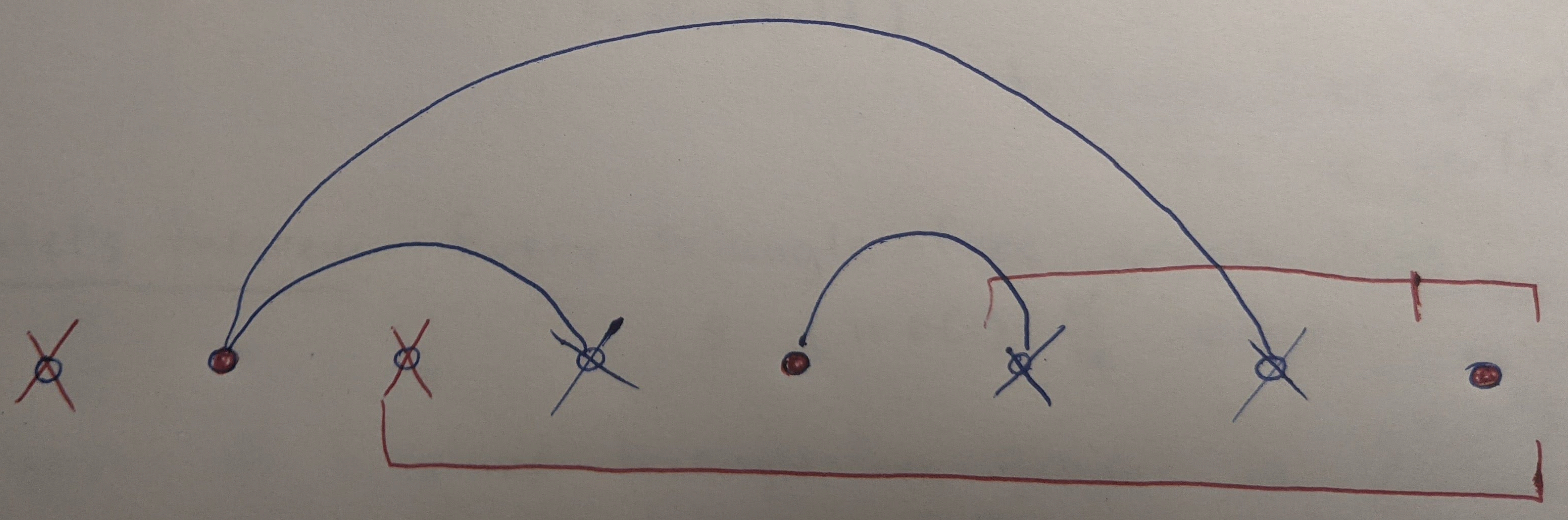


Lecture 22:

Hypergraph Containers



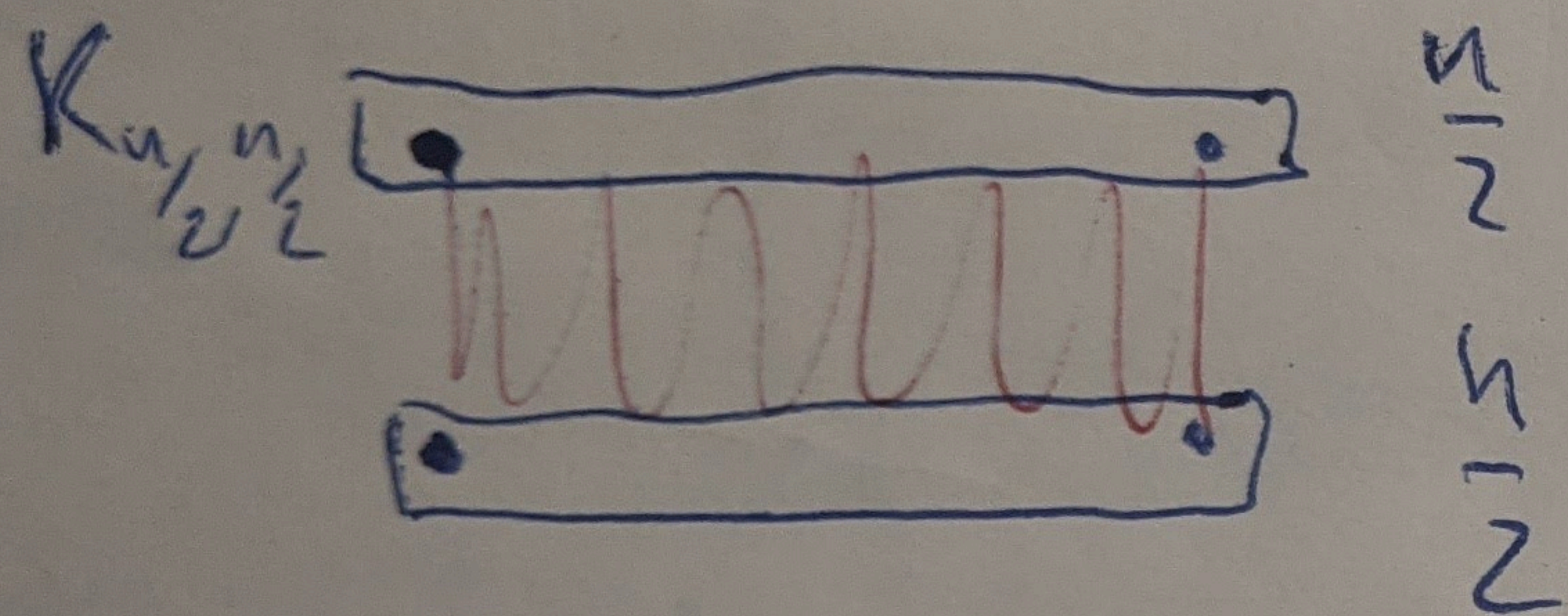
12. Hypergraph Containers

(Balogh, Morris, Samotij & Saxton, Thomason, 2015)

"Good" approximations of independent sets in Hypergraphs.

How many triangle-free graphs are there on n vertices?

Lower bound:



$\frac{n^2}{4}$ edges $\rightarrow 2^{\frac{n^2}{4}}$ subgraphs.

Theorem
(Erdős, Kleitman, Rothschild)

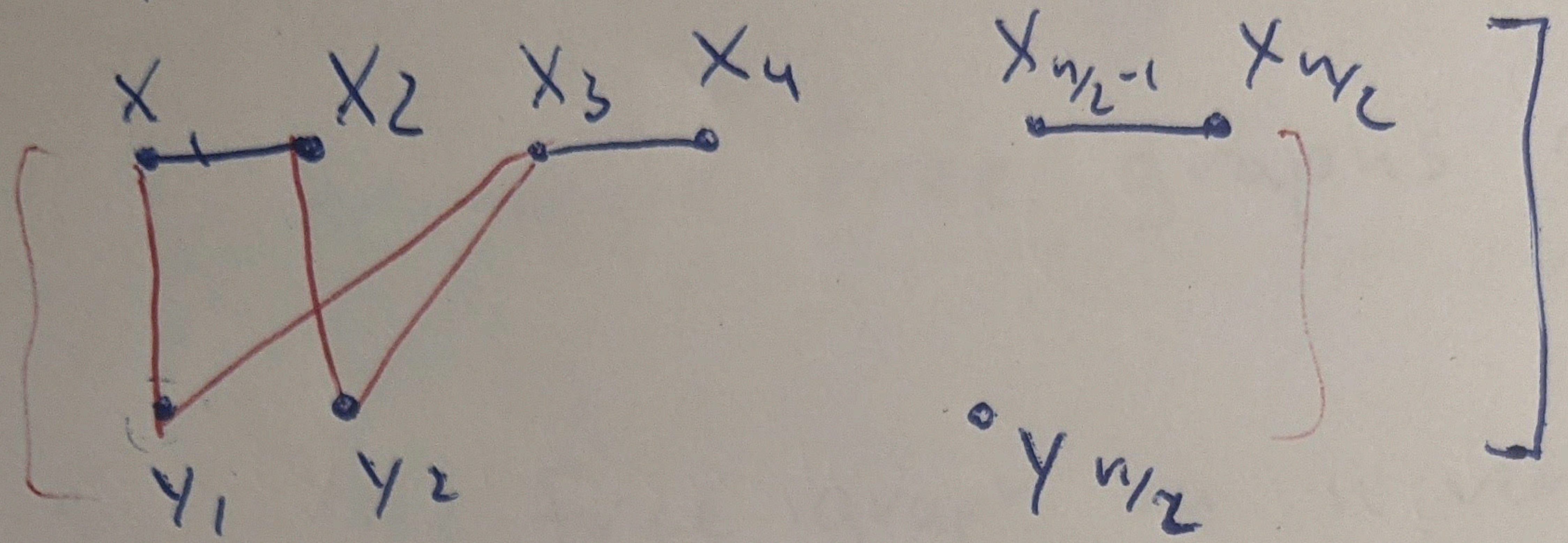
There are $(1+o(1)) 2^{\frac{n^2}{4}}$ triangle-free graphs on n vertices.

Mantel's theorem: Every triangle-free graph has $\leq (1+o(1)) \frac{n^2}{4}$ edges on n vertices.

If number of maximal triangle-free graphs is $2^{o(n^2)}$ then the theorem follows from Mantel by taking union.

Unfortunately, it's not true:

Example:



join each y_i to exactly one of vertices

x_{2j-1}, x_{2j} for every j

There are

2

maximal triangle free graphs of this form.

Subgraphs of a universal graph which is not triangle-free but still has $(1+o(1)) \frac{n^2}{4}$ edges (and is "close" to being triangle-free).

Container

Theorem 12.1: For every $\epsilon > 0$ there exist $C > 0$. Such that for every n there exists a collection

\mathcal{C}

of n vertex graphs such that.

containers for triangle free graphs.

- every triangle free graph on n vertices is a subgraph of a graph in \mathcal{C}
- every graph in \mathcal{C} has at most $(\frac{1}{4} + \epsilon) n^2$ edges
- $|\mathcal{C}| \leq C n^{3/2}$

Corollary 12.2:

EKR

There are

~~$2^{(1+o(1))n^2}$~~ $2^{(1+o(1))\frac{n^2}{4}}$

triangle free graphs on n vertices

Proof:

Every triangle free graph on n vertices is a subgraphⁿ of some graph in \mathcal{C} , and each $G \in \mathcal{C}$ has

$\leq 2^{(\frac{1}{4} + \epsilon)n^2}$ subgraphs

By union bound there are at most

$2^{(\frac{1}{4} + \epsilon)n^2} |\mathcal{C}| \leq 2^{(\frac{1}{4} + \epsilon)n^2 + cn^{3/2} \log n}$
triangle free graphs.

Corollary 12.3:

For $p \gg \frac{\log n}{\sqrt{n}}$ with high probability every triangle free subgraph of a graph in $G(n, p)$

↓
Mantel's
theorem
for random
graphs.

has $\leq (\frac{1}{4} + o(1))pn^2$ edges.

Proof: Let $G \in \mathcal{L}$
 By Chernoff bound

$$\mathbb{P} \left[\underbrace{|E(G) \cap E(G(n,p))|}_{\substack{\downarrow \\ \text{expected size} \\ \leftarrow p |E(G)|}} \geq (\rho + \varepsilon) |E(G)| \right] \leq e^{-c_\varepsilon n^2 p}$$

$(\frac{p}{4} + 3\varepsilon)n^2$
 $(\frac{1}{4} + \varepsilon)n^2$

Taking union over $G \in \mathcal{L}$

$$\mathbb{P} \left[G(n,p) \text{ contains a triangle-free subgraph with } \geq \left(\frac{p}{4} + 3\varepsilon \right) n^2 \text{ edges} \right] \leq e^{-c_\varepsilon n^2 p} \cdot e^{cn^{3/2} \log n}$$

$(1+3\varepsilon) \frac{pn^2}{4}$
 \neq instead

$$n^2 p \gg n^{3/2} \log n$$

$$p \gg \frac{\log n}{\sqrt{n}} \quad \checkmark$$

Construct a 3-uniform hypergraph \mathcal{H}
by setting

$$V(\mathcal{H}) = [n]^{(2)}$$

↓ set of all edges of ~~complete~~
 K_n

$$E(\mathcal{H}) = \{ e_1, e_2, e_3, \dots, e_1, e_2, e_3 \text{ form a triangle} \}$$

Graph on n vertices
triangle free

↔ subsets of $V(\mathcal{H})$

↔ subset of $V(\mathcal{H})$ is

independent

does not contain all vertices
of any edge.

Theorem 12.1 follows from a more general
container theorem for independent
sets in 3-uniform hypergraphs.

Theorem 12.4: For every $c > 0$ there exist $\delta > 0$,
satisfying the following.

Let H be a 3-uniform hypergraph
with average degree $\underline{d} \geq \frac{1}{\delta}$
s.t. no vertex has degree $> c \cdot d$
no pair of vertices belongs
to more than $c\sqrt{d}$ edges together.

for initial H
for 12.1
 $\underline{d} = \frac{n-1}{2}$

Then there exists \mathcal{C} collection of subsets of $V(H)$

in triangle-free graphs on n vertices. s.t.

- every independent set in H is contained in some $C \in \mathcal{C}$

$(1-\delta) \binom{n}{2} \rightarrow$ want $(\frac{1}{4} + \epsilon)n^2$

$|C| \leq (1-\delta) |V(H)|$

$|\mathcal{C}| \leq n$

where $n = |V(H)|$

if a container C
has $> (\frac{1}{4} + \epsilon)n^2$ edges
then ϵn^3 triangles.

apply theorem 12.4
again.

Theorem 12.5: For every $\epsilon > 0$ there exists $\delta > 0$ satisfying the following

Let G be a graph with average degree d and maximum degree $\leq Cd$

Then there exists \mathcal{C} collection of subsets of $V(G)$ s.t.

- every ind. set of G is contained in some $C \in \mathcal{C}$

- $|C| \leq (1-\delta)|V(G)|$.

- $|\mathcal{C}| \leq n^{\frac{2\delta n}{d}}$ where $n = |V(G)|$.

Proof: ~~Algorithm. Input: G~~

We will show that for every independent set I in G there exists a Signature $S(I) \subseteq I$

s.t. • $S(I) \subseteq I \subseteq C(S(I))$

container for I

there are few signatures

determined only by its signature $S(I)$.

$|C(I)| \leq \frac{2\delta n}{d}$

$|C(S(I))| \leq (1-\delta)n$

Algorithm for finding $S(I)$ and $C(S(I))$.

Input: I - independent.

Maintain:
[S - current signature
 A - set of available vertices
 X - set of forbidden vertices]

Start:

$$S = \emptyset$$

$$A = V(G)$$

$$X = \emptyset$$

Steps: While $|S| \leq \frac{2\delta n}{d}$

Let $v \in A \cap I$

with maximum degree in $G[A]$
(where ties are broken in some canonical order)

$$- S := S \cup \{v\}$$

$$- X := X \cup \{ \text{neighbors of } v \} \cup \{ \text{vertices of } A \text{ with higher degree than } v \text{ in } G[A] \}$$

$$- A := V(G) - S - X.$$

Output: $S(I) = S$, $C(I) = V(G) - X$.
" " $C(S(I))$

~~Key~~

Need: $|C(S(I))| = |V(G) - X| \leq (1 - \delta)n$ if $|X| \geq \delta n$

\rightarrow [$X \cap I = \emptyset$ throughout.
 $(S; A, X)$ is a partition of $V(G)$
 $S \subseteq I$]

We stop when $|S| = \frac{2\delta n}{d}$ we made $\frac{2\delta n}{d}$ step.

Enough to show that X is increased by $\geq \frac{d}{2}$ at each step.

assuming for a contradiction $|X| \leq \delta n$.

Suppose ~~on~~ Consider the first step when X is increased by $< \frac{d}{2}$.

maximum degree

So far we have deleted $\leq \left(\frac{2\delta n}{d} + \delta n \right) cd$ edges
 for s vertices

then at most $\frac{d}{2}$ vertices precede the last chosen v in its order in $G[A]$

$$|E(G)| \leq \left(\frac{2\delta n}{d} + \delta n \right) cd$$

$$+ \underbrace{cd^2}_{\frac{d}{2}} + \underbrace{\frac{d}{2}n}_{\frac{d}{2}n} < \frac{d}{2}n = |E(G)| \text{ and } \deg(v) \leq \frac{d}{2}$$

if δ is small enough

$$2E(G[A]) \leq \underbrace{cd \cdot \frac{d}{2}}_{\text{degrees before } v} + \underbrace{\frac{d}{2} \cdot n}_{\text{after}}$$